## Problem A

Asynchronous Exceptions

Time Limit: 5 sec , Memory Limit: 64 MB

Your job is to simulate the operation on a machine called ACM (Asynchrounous Calculating Machine) to determine how long each program takes to finish their job. This helps detecting unexpected thread behaviors, like race and deadlock.

ACM has many processing units and can run many threads simultaneously. It runs a global scheduler, which maintains all the threads, assigns a thread on each processing unit and manages all semaphores. Each thread has its own thread id and every thread is one of three states: RUNNING, READY, and WAITING. Global scheduler has one thread queue which contains all the threads in READY state.

## explanation of a program

A program consists of one or more code blocks. Each code block has a list of ACM operations, including the actual computations and thread creation. Each thread runs operations in only one of the given code blocks, and it runs operations from the top of the block to the bottom, unless it receives killThread-signal from other threads.

## explanation of operations

ACM simulation program simulates 9 operations. Here I describe them all. Please note that word enclosed by [] should have some actual name or value.
compute [clock] Spend [clock] clocks for computation. When this operation is interrupted after spending T clocks by other threads or the scheduler and is resumed later, it consumes [clock] - T clocks.
[thread-var] <- forkR [code block] Generate a native thread running on [code block] and store its id in a thread variable [thread-var]. It can run simultaneously with any threads which are generated in other processing unit. If [thread-var] is used before in the thread, its previous value is overwritten and lost. (This means that we will not be able to refer to the thread although it might run forever.)
[thread-var] <- forkI [code block] Generate a virtual thread running on [code block] and store its id in a thread variable [thread-var]. A virtual thread shares some OS resources with its parent thread, so they can't run simultaneously even there are idle processing units. This restriction is applied to any two threads which are directly or indirectly connected by the parent-child relatinship of forkI. If [thread-var] is used before in the thread, its previous value is overwritten and lost.
yield Make the current thread from running state to ready state. The thread goes to the end of the thread queue.
killThread [thread-var] Send a kill signal to the [thread-var] thread. [thread-var] guaranteed to be used in either forkR or forkI operations in the same code block before this operation. When a thread receives a kill signal, it immediately cancels its resource requests by lock, changes into ready state if it's in block state, and ends its operation when it's in running state. This operation does not change the semaphore values. If the target thread is already ended, this operation does nothing.
lock [semaphore name] [amount] Request [amount] of [semaphore name]. If the value of [semaphore name] is greater or equal to the [amount], it just subtracts the [amount] from the semaphore variable. Otherwise, the operation blocks the current thread until the variable is greater or equal to the [amount]. Once the variable gets greater or equal to the [amount], the thread gets into READY state and goes into the thread queue after subtracting [amount] from the variable. If there are more than one thread requesting for the same semaphore, the thread which locked the semaphore first
always gets READY first. So, the other threads won't get READY even if the semaphore variable meets their demands. If there are more than one thread getting READY, the thread which locked the semaphore earlier goes into the queue first.
unlock [semaphore name] [amount] Add [amount] to [semaphore name]. Values of semaphores can be larger than initial values. The thread doesn't get blocked at all by this operation.
loop [loop count] Run the code snippet between loop and corresponding next command for [loop count] times.
next This corresponds a loop command in the same code block. The code snippet between loop and next should be run.
end End the operation. This command only appears in the end of the code block, and each code block has this operation at the end.

Each parameter should either be a name or a digit. A name is an alphabetic string (case-sensitive) and its length is at most 200. Here are some explanations.
[thread-var] It is a name. It only refers to the id in the same thread. Each thread has its own value even it has the same name.
[clock] It is a non-negative integer which is at most 1,000 .
[semaphore-name] It is a name.
[amount] It is a non-negative integer which is at most 1,000 .
[loop count] It is a non-negative integer which is at most 1,000 .

Additionally, you may assume following constraints.

- Total number of lines that all threads evaluate throughout simulation never exceeds 100,000 .
- Size of input file never exceeds 100 KB .


## explanation of a simulator

Threads are assigned to CPUs when

- a simulation starts. (A thread is assigned to CPU 1. Its code block is the first code block of the input.)
- there are more than one threads that are in ready state which can be assigned to a CPU. (Threads are asssinged in the thread queue order. If there are more than one CPUs that can be assinged to the thread, the smallest CPU is selected.)
- time step is a multiple of time slice. (If the CPU has a thread that is running state, the thread goes to ready state in asscending order of the CPU id. After this, threads are assigned to each CPU in asscending order of the CPU id.)

In each step, the simulator does round-robin preemption (when time step is a multiple of time slice), executes operations of each running threads and increments time step. When a compute operation is executed, the thread gets computing time. A thread that has computing time greater than 0 cannot do any operations. After each step, the simulator decrements positive computing time of each running thread.

Operations executed in the same time step is executed in ascending order of the CPU id. In the same time step, a CPU id of an operation is greater or equal to CPU ids of operations which have already been executed.

Time step starts from 0 .

## Input

Input consists of multiple testcases.
Each testcase begins with a line that contains two integers separated by a space. They mean the number of time steps of the simulation and the maximum number of threads that ACM is capable, respectively. The next line contains a single integer, which indicates the number of CPUs available. The following line has an integer that means time slice of the scheduler.

Description of semaphores follows. It begins with the number of semaphores which is followed by information of each semaphore, one per line. The information of a semaphore consists of an alphabet string (case-sensitive) and an integer, which specify the name of the semaphore and its initial value, respectively.

Finally code blocks are given. The number of code blocks comes first, and then each code block is described. The first line of the description of each code block is the name of the block (case-sensitive alphabet string), followed by a colon (:). Following lines describe content of the code block. A code block is an array of operations described above, and one operation is written in one line. You may assume that every code block always ends with an 'end' operation.

Input terminates with two zeroes separated by a space character.

## constraints

- Threre is 1 thread queue.
- The id of the $n$-th created thread is $n$.
- CPU ids are 1 though the number of CPUs.
- The number of time steps is at most 1,000 .
- The maximum number of threads is at most 1,000 .
- The number of CPUs is at most 100 .
- Time slice is at most 1,000
- The number of semaphores is at most 1,000 .
- The number of code blocks is at most 1,000 .


## Output

For each test case, print its case number on the first line.
If the number of living threads exceeds ACM's capacity, put information of all threads which terminated before exceeding the capacity, and then put "<<oops>>". If not, and if there are one or more threads not finished at the end of the simulation, put information of all threads terminated, and then put "<<loop>>". Otherwise (i.e. when all threads terminates within the time), just put information of all threads. All quotes are for clarity.

Information of threads must be written in increasing order of thread ID, one per line. Information of a single thread is denoted by two integers, the thread ID and the time it terminated, separated by a single space character.

## Sample Input

```
50 50
1
10
1
semaphore 1
1
codeBlockA:
loop 2
compute 10
next
end
50 50
1
10
1
semaphore 1
2
codeBlockA:
hoge <- forkR codeBlockB
yield
compute 1
killThread hoge
lock semaphore 1
compute 1
end
codeBlockB:
compute 1
lock semaphore 2
end
55
1
3
1
semaphore 1
1
codeBlockA:
hoge <- forkI codeBlockA
compute 1
end
5 5
1
2
1
semaphore 1
1
codeBlockA:
compute 1
hoge <- forkI codeBlockA
hoge <- forkI codeBlockA
end
0
```


## Output for the Sample Input

```
Case 1:
120
Case 2:
1 3
2 3
Case 3:
11
2
34
44
5 5
<<loop>>
Case 4:
11
2 3
3 
<<oops>>
```


## Problem B AYBABTU

There is a tree that has $n$ nodes and $n-1$ edges. There are military bases on $t$ out of the $n$ nodes. We want to disconnect the bases as much as possible by destroying $k$ edges. The tree will be split into $k+1$ regions when we destroy $k$ edges. Given the purpose to disconnect the bases, we only consider to split in a way that each of these $k+1$ regions has at least one base. When we destroy an edge, we must pay destroying cost. Find the minimum destroying cost to split the tree.

## Input

The input consists of multiple data sets. Each data set has the following format. The first line consists of three integers $n, t$, and $k(1 \leq n \leq 10,000,1 \leq t \leq n, 0 \leq k \leq t-1)$. Each of the next $n-1$ lines consists of three integers representing an edge. The first two integers represent node numbers connected by the edge. A node number is a positive integer less than or equal to $n$. The last one integer represents destroying cost. Destroying cost is a non-negative integer less than or equal to 10,000 . The next $t$ lines contain a distinct list of integers one in each line, and represent the list of nodes with bases. The input ends with a line containing three zeros, which should not be processed.

## Output

For each test case, print its case number and the minimum destroying cost to split the tree with the case number.

Sample Input

```
2 1
21
1
2
3 2
2 1
132
143
2
3
4
000
```

Output for the Sample Input

```
Case 1: 1
Case 2: 3
```


# Problem C <br> Billiards Sorting 

Time Limit: 7 sec , Memory Limit: 64 MB

Rotation is one of several popular pocket billiards games. It uses 15 balls numbered from 1 to 15 , and set them up as illustrated in the following figure at the beginning of a game. (Note: the ball order is modified from real-world Rotation rules for simplicity of the problem.)

```
            [ 1]
            [ 2][ 3]
    [ 4][ 5][ 6]
    [ 7][ 8][ 9][10]
[11] [12] [13] [14] [15]
```

You are an engineer developing an automatic billiards machine. For the first step you had to build a machine that sets up the initial condition. This project went well, and finally made up a machine that could arrange the balls in the triangular shape. However, unfortunately, it could not place the balls in the correct order.

So now you are trying to build another machine that fixes the order of balls by swapping them. To cut off the cost, it is only allowed to swap the ball \#1 with its neighboring balls that are not in the same row. For example, in the case below, only the following pairs can be swapped: $(1,2),(1,3),(1,8)$, and $(1,9)$.

```
            [ 5]
            [ 2][ 3]
            [ 4][ 1][ 6]
    [ 7][ 8][ 9][10]
[11][12][13][14][15]
```

Write a program that calculates the minimum number of swaps required.

## Input

The first line of each test case has an integer $N(1 \leq N \leq 5)$, which is the number of the rows.
The following $N$ lines describe how the balls are arranged by the first machine; the $i$-th of them consists of exactly $i$ integers, which are the ball numbers.

The input terminates when $N=0$. Your program must not output for this case.

## Output

For each test case, print its case number and the minimum number of swaps.
You can assume that any arrangement can be fixed by not more than 45 swaps.

## Sample Input

```
2
3
2 1
4
9
24
8 5 3
71610
0
```

Output for the Sample Input
Case 1: 1
Case 2: 13

## Problem D <br> Digit

For a positive integer $a$, let $S(a)$ be the sum of the digits in base $l$. Also let $L(a)$ be the minimum $k$ such that $S^{k}(a)$ is less than or equal to $l-1$. Find the minimum $a$ such that $L(a)=N$ for a given $N$, and print $a$ modulo $m$.

## Input

The input contains several test cases, followed by a line containing " 000 ". Each test case is given by a line with three integers $N, m, l\left(0 \leq N \leq 10^{5}, 1 \leq m \leq 10^{9}, 2 \leq l \leq 10^{9}\right)$.

## Output

For each test case, print its case number and the minimum $a$ modulo $m$ as described above.

## Sample Input

```
0 1000 10
1 1000 10
0 0
```


## Output for the Sample Input

```
Case 1: 1
Case 2: 10
```


## Problem E <br> Dungeon Creation

The king demon is waiting in his dungeon to defeat a brave man. His dungeon consists of $H \times W$ grids. Each cell is connected to four (i.e. north, south, east and west) neighboring cells and some cells are occupied by obstacles.

To attack the brave man, the king demon created and sent a servant that walks around in the dungeon. However, soon the king demon found that the servant does not work as he wants. The servant is too dumb. If the dungeon had cyclic path, it might keep walking along the cycle forever.

In order to make sure that the servant eventually find the brave man, the king demon decided to eliminate all cycles by building walls between cells. At the same time, he must be careful so that there is at least one path between any two cells not occupied by obstacles.

Your task is to write a program that computes in how many ways the kind demon can build walls.

## Input

The first line of each test case has two integers $H$ and $W(1 \leq H \leq 500,1 \leq W \leq 15)$, which are the height and the width of the dungeon. Following $H$ lines consist of exactly $W$ letters each of which is '.' (there is no obstacle on the cell) or ' $\#$ ' (there is an obstacle). You may assume that there is at least one cell that does not have an obstacle.

The input terminates when $H=0$ and $W=0$. Your program must not output for this case.

## Output

For each test case, print its case number and the number of ways that walls can be built in one line. Since the answer can be very big, output in modulo $1,000,000,007$.

## Sample Input

```
22
..
..
3 3
...
...
..#
0
```


## Output for the Sample Input

```
Case 1: 4
Case 2: 56
```


## Problem F Longest Lane

Mr. KM, the mayor of KM city, decided to build a new elementary school. The site for the school has an awkward polygonal shape, which caused several problems. The most serious problem was that there was not enough space for a short distance racetrack. Your task is to help Mr. KM to calculate the maximum possible length for the racetrack that can be built in the site. The track can be considered as a straight line segment whose width can be ignored. The boundary of the site has a simple polygonal shape without self-intersection, and the track can touch the boundary. Note that the boundary might not be convex.

## Input

The input consists of multiple test cases, followed by a line containing " 0 ". Each test case has the following format. The first line contains an integer $N(3 \leq N \leq 100)$. Each of the following $N$ lines contains two integers $x_{i}$ and $y_{i}\left(-1,000 \leq x_{i}, y_{i} \leq 1,000\right)$, which describe the coordinates of a vertex of the polygonal border of the site, in counterclockwise order.

## Output

For each test case, print its case number and the maximum possible length of the track in a line. The answer should be given as a floating point number with an absolute error of at most $10^{-6}$.

## Sample Input

```
4
0}
10 0
10 10
010
3
0
10
0 1
0
```


## Output for the Sample Input

```
Case 1: 14.142135624
Case 2: 1.41421356
```


## Problem G <br> PLAY in BASIC

Time Limit: 5 sec , Memory Limit: 64 MB
15 Apr 2012

One sunny day, Dick T. Mustang found an ancient personal computer in the closet. He brought it back to his room and turned it on with a sense of nostalgia, seeing a message coming on the screen:

## READY?

Yes. BASIC.
BASIC is a programming language designed for beginners, and was widely used from 1980 's to 1990's. There have been quite a few BASIC dialects. Some of them provide the PLAY statement, which plays music when called with a string of a musical score written in Music Macro Language (MML). Dick found this statement is available on his computer and accepts the following commands in MML:

Notes: Cn, C+n, C-n, Dn, D+n, D-n, En,....... ( $\mathbf{n}=\mathbf{1 , 2 , 4 , 8 , 1 6 , 3 2 , 6 4 , 1 2 8}+$ dots) Each note command consists of a musical note followed by a duration specifier.
Each musical note is one of the seven basic notes: ' $C$ ', ' $D$ ', ' $E$ ', ' $F$ ', ' $G$ ', ' $A$ ', and ' $B$ '. It can be followed by either ' + ' (indicating a sharp) or '-' (a flat). The notes ' C ' through ' B ' form an octave as depicted in the figure below. The octave for each command is determined by the current octave, which is set by the octave commands as described later. It is not possible to play the note ' C -' of the lowest octave (1) nor the note ' $\mathrm{B}+$ ' of the highest octave (8).
Each duration specifier is basically one of the following numbers: ' 1 ', ' 2 ', ' 4 ', ' 8 ', ' 16 ', ' 32 ', ' 64 ', and ' 128 ', where ' 1 ' denotes a whole note, ' 2 ' a half note, ' 4 ' a quarter note, ' 8 ' an eighth note, and so on. This specifier is optional; when omitted, the duration will be the default one set by the L command (which is described below). In addition, duration specifiers can contain dots next to the numbers. A dot adds the half duration of the basic note. For example, '4.' denotes the duration of ' 4 ' (a quarter) plus ' 8 ' (an eighth, i.e. half of a quarter), or as 1.5 times long as ' 4 '. It is possible that a single note has more than one dot, where each extra dot extends the duration by half of the previous one. For example, ' $4 .$. ' denotes the duration of ' 4 ' plus ' 8 ' plus ' 16 ', ' $4 . .$. ' denotes the duration of ' 4 ' plus ' 8 ' plus ' 16 ' plus ' 32 ', and so on. The duration extended by dots cannot be shorter than that of ' 128 ' due to limitation of Dick's computer; therefore neither ' 128 .' nor ' $32 \ldots$ '.' will be accepted. The dots specified without the numbers extend the default duration. For example, ' C .' is equivalent to ' C 4 .' when the default duration is ' 4 '. Note that ' C 4 C 8 ' and 'C4.' are unequivalent; the former contains two distinct notes, while the latter just one note.

Rest: $\mathbf{R n}(\mathbf{n}=\mathbf{1 , 2 , 4 , 8 , 1 6 , 3 2 , 6 4 , 1 2 8}+$ dots) The $R$ command rests for the specified duration. The duration should be specified in the same way as the note commands, and it can be omitted, too. Note that 'R4R8' and 'R4.' are equivalent, unlike 'C4C8' and 'C4.', since the both rest for the same duration of ' 4 ' plus ' 8 '.

Octave: On $(\mathbf{n}=\mathbf{1 - 8}),<,>$ The O command sets the current octave to the specified number. ' $>$ ' raises one octave up, and ' $<$ ' drops one down. It is not allowed to set the octave beyond the range from 1 to 8 by these commands. The octave is initially set to 4 .

Default duration: $\operatorname{Ln}(\mathbf{n}=\mathbf{1 , 2 , 4 , 8 , 1 6 , 3 2}, \mathbf{6 4}, 128) ~ T h e ~ L ~ c o m m a n d ~ s e t s ~ t h e ~ d e f a u l t ~ d u r a t i o n . ~ T h e ~ d u r a-~$ tion should be specified in the same way as the note commands, but cannot be omitted nor followed by dots. The default duration is initially set to 4 .

Volume: Vn ( $\mathbf{n}=\mathbf{1 - 2 5 5})$ The V command sets the current volume. Larger is louder. The volume is initially set to 100 .


As an amateur composer, Dick decided to play his pieces of music by the PLAY statement. He managed to write a program with a long MML sequence, and attempted to run the program to play his music but unfortunately he encountered an unexpected error: the MML sequence was too long to be handled in his computer's small memory!

Since he didn't want to give up all the efforts he had made to use the PLAY statement, he decided immediately to make the MML sequence shorter so that it fits in the small memory. It was too hard for him, though. So he asked you to write a program that, for each given MML sequence, prints the shortest MML sequence (i.e. MML sequence containing minimum number of characters) that expresses the same music as the given one. Note that the final values of octave, volume, and default duration can be differrent from the original MML sequence.

## Input

The input consists of multiple data sets. Each data set is given by a single line that contains an MML sequence up to 100,000 characters. All sequences only contain the commands described above. Note that each sequence contains at least one note, and there is no rest before the first note and after the last note.

The end of input is indicated by a line that only contains "**". This is not part of any data sets and hence should not be processed.

## Output

For each test case, print its case number and the shortest MML sequence on a line.
If there are multiple solutions, print any of them.
Sample Input

```
C4C4G4G4A4A4G2F4F4E4E4D4D4C2
04C4.C8F4.F8G8F8E8D8C2
B-8>C8<B-8V40R2R . .R8.V100EV50L1CG
*
```


## Output for the Sample Input

[^0]
# Problem H Skyland 

Somewhere in the sky, KM kingdom built $n$ floating islands by their highly developed technology. The islands are numbered from 1 to $n$.

The king of the country, Kita_masa, can choose any non-negative real number as the altitude for each island, as long as the sum of the altitudes is greater than or equals to $H$. For floating the island $i$ to the altitude $h_{i}$, it costs $b_{i} h_{i}$. Besides, it costs $\left|h_{i}-h_{j}\right| c_{i, j}$ for each pair of islands $i$ and $j$ since there are communications between these islands.

Recently, energy prices are rising, so the king, Kita_masa, wants to minimize the sum of the costs. The king ordered you, a court programmer, to write a program that finds the altitudes of the floating islands that minimize the cost.

## Input

The input contains several test cases. Each test case starts with a line containing two integers $n(1 \leq n \leq$ $100)$ and $H(0 \leq H \leq 1,000)$, separated by a single space. The next line contains $n$ integers $b_{1}, b_{2}, \ldots, b_{n}$ ( $0 \leq b_{i} \leq 1,000$ ). Each of the next $n$ lines contains $n$ integers $c_{i, j}\left(0 \leq c_{i, j} \leq 1,000\right)$. You may assume $c_{i, i}=0$ and $c_{i, j}=c_{j, i}$.

The last test case is followed by a line containing two zeros.

## Output

For each test case, print its case number. Then print a line containing a space-separated list of the altitudes of the islands that minimizes the sum of the costs. If there are several possible solutions, print any of them. Your answer will be accepted if the altitude of each island is non-negative, sum of the altitudes is greater than $\left(1-10^{-9}\right) \mathrm{H}$, and the cost calculated from your answer has an absolute or relative error less than $10^{-9}$ from the optimal solution.

Follow the format of the sample output.

## Sample Input

```
2 1
13
01
10
3 3
124
020
2 01
010
00
```

Output for the Sample Input

```
Case 1:
0.750 .25
Case 2:
1.51 .50 .0
```


# Problem I <br> Three-way Branch 

Time Limit: 7 sec , Memory Limit: 64 MB

There is a grid that consists of $W \times H$ cells. The upper-left-most cell is $(1,1)$. You are standing on the cell of $(1,1)$ and you are going to move to cell of $(W, H)$. You can only move to adjacent lower-left, lower or lower-right cells.

There are obstructions on several cells. You can not move to it. You cannot move out the grid, either. Write a program that outputs the number of ways to reach $(W, H)$ modulo $1,000,000,009$. You can assume that there is no obstruction at $(1,1)$.

## Input

The first line contains three integers, the width $W$, the height $H$, and the number of obstructions $N$. $\left(1 \leq W \leq 75,2 \leq H \leq 10^{18}, 0 \leq N \leq 30\right)$ Each of following $N$ lines contains 2 integers, denoting the position of an obstruction $\left(x_{i}, y_{i}\right)$.

The last test case is followed by a line containing three zeros.

## Output

For each test case, print its case number and the number of ways to reach $(W, H)$ modulo $1,000,000,009$.

## Sample Input

```
241
2 1
2 1
2}
0 0
```

Output for the Sample Input

```
Case 1: 4
Case 2: 0
```


## Problem J Tree Allocation

A tree is one of the most popular data structures in computer science. Tree itself is very elegant, but it is incompatible with current computer architecture. In the current architecture, memory can be regarded as a one-dimensional array. Since a tree is not a one-dimensional structure, cache efficiency may come into question when we allocate a tree on memory. We can access the data faster if it is contained in cache. Let us consider the allocation of nodes of a tree which achieves high cache performance.

We define the cost of allocation for a tree as follows. For simplicity, we regard memory to be separated to blocks each can store at most $B$ nodes of a tree. When we access data in a block, all data in the block are stored in cache and access request to data in the block will be faster. The cost of an access to node $u$ after node $v$ is 0 if $u$ and $v$ are in same block, and 1 otherwise. Initially, cache has no data and the cost of an access to a node is always 1 . The cost of the path $v_{1}, v_{2}, \ldots, v_{n}$ is sum of the cost when we access the nodes $v_{1}, v_{2}, \ldots, v_{n}$ in this order. For a tree, we define the cost of allocation as a maximum cost of the paths from root node to each terminal node.

The figures below show examples of allocation when $B=4$ and node 1 is the root node (it corresponds to the tree described in the third sample input). Each frame represents a block. The left figure is an example of allocation whose cost is 3 . It is not optimal, because the right example achieves the optimal allocation cost 2.


Given a tree, for each node $i$ in the tree, calculate the minimum cost of allocation when the node $i$ is the root node.

## Input

The input contains several test cases. Each test case starts with a line containing two integers $N(1 \leq$ $N \leq 100,000)$ and $B(1 \leq B \leq N)$, separated by a single space. Each of the next $N-1$ lines contains an integer. The $i$-th integer $p_{i}\left(1 \leq p_{i} \leq i\right)$ denotes that the node $i+1$ and the node $p_{i}$ are connected. The nodes are numbered from 1 to $N$.

The last test case is followed by a line containing two zeros.

## Output

For each test case, print its case number. Then print $N$ lines. The $i$-th line should contain an integer which denotes the minimum cost of allocation of the given tree when node $i$ is the root node.

Follow the format of the sample output.

## Sample Input

```
3 1
1
2
3 2
1
1
O
1
1
2
3
3
4
4
4
5
0
```


## Output for the Sample Input

```
Case 1:
```

3
2
3
Case 2:
2
2
2
Case 3:
2
2
2


[^0]:    Case 1: CCGGAAG2FFEEDDC2
    Case 2: C.C8F.L8FGFEDC2
    Case 3: L8B-B+B-RL1RE4V50CG

